LECTURE 24

Theory and Design of PL (CS 538) April 20, 2020



HW5 WRAPUP

An enormous pain in the ass
Why go through all of this trouble?
Ownership: rule out most memory leaks
Aliasing: make totally unsafe stuff safe

BUILD A TREE THAT...

- With no runtime checks (direct memory write) While never segfaulting or breaking the tree
- Gives clients pointers to internal tree memory • Lets client write whenever and whatever they want • Also build an iterator handing out pointers all over



HW5: FEEDBACK?



HW6 OUT

- Concurrency: making things go faster
- We give you: a slow, single-threaded version
- You make: multi-threaded versions in two ways
- Should be much less grappling with borrow checker
 But still a bit (The Rust Programming Language)

Start early, especially if you haven't tried writing concurrent code before

gs go faster e-threaded version versions in two ways oling with borrow checker Programming Language)

MODELING Concurrency

MANY ASPECTS

- Parallelism and simultaneous execution
- Message-passing and channels
- Shared memory and locking
- Threads blocking/waiting

Itaneous execution d channels locking aiting

PROCESS CALCULUS

• Mathematical model of message passing Many flavors, developed in 1970s and 1980s Communicating sequential processes (CSP) Communicating concurrent systems (CCS) Pi-calculus

• By Tony Hoare, Robin Milner, and many others

OUR VERSION

Simple arithemtic expressions
 Channels: named pipes for communication
 Processes: send/receive along channels

PROCESS CALCULUS: GRANNAR



ARITHMETIC EXPRESSIONS

num	=	" () "		"1"	"2"	• • •		
var	=	"×"		"Y"	" Z "	•••		
exp	=	num		var	exp	"+" exp		exp

• Arithmetic expressions with variables • Examples of arithmetic expressions: **4**2 **5***3 ■ 2 + 0 + z

p "*" exp | ...

CHANNELS

chn = "A" | "B" | "C" | ...

- Addresses to send to/receive from
- Different names = different addresses
- We'll use two special channels:
 - I: input channel into program
 - O: output channel from program

PROCESSES

prc =	"done"	(*	do nothing	*)
L 1	"make" chn "in" prc	(*	make new channel	*)
L 1	"send" exp "->" chn "then" prc	(*	send a message	*)
L 1	"recv" var "<-" chn "then" prc	(*	receive a message	*)
L 1	"[" exp "<" exp "]" prc	(*	run if guard true	*)
L 1	prc "+" prc	(*	run this or that	*)
	prc " " prc	(*	run in parallel	*)

- Make new channel, send and receive along channel
- Combine several processes together
 Select between different processes
 - Run processes in parallel

and receive along channel es together at processes el

EXAMPLES (BLACKBOARD)

OPERATIONAL SEMANTICS



MAIN SETUP

• Define how processes step $\mathrm{P} \to \mathrm{Q}$ • New addition: each transition may have a label Labels model sending and receiving (A, n): send num n along channel A • (\overline{A}, n) : receive num n from channel A • Other steps: no label (silent transitions)

BLACKBOARD (OR WR6)



EXTENSION: RECURSION

WHY RECURSION?

- So far: finite number of steps
- Some processes live forever (e.g., servers)
- Extend the language with recursive processes

f steps prever (e.g., servers) /ith *recursive processes*

Add process names and recursive definitions

name = "P1" "P2" "P3"	(
prc = name	(
def = name "=" prc	(





* names of processes *)

* process could be a name *)

(* definition of processes *)



OPERATIONAL SEMANTICS Just add one more rule to unfold definitions...

A TINY GLIMPSE OF



JOE ARMSTRONG Invented Erlang while working for Ericsson • Hugely influential views on computing

- Passed away in 2019 :(
- - Take a look at his thesis
 - Or check out some of his talks

PRINCIPLE 1: PROCESSES

- Take idea of process from OS
 Not threads: no shared memory space!
 - Separate program into several processes
- Erlang: processes are cheap
 - Can make millions of processes
 - So-called "green threads"
- Rust: heavier, OS threads (can't have so many)
 Used to have green threads, taken out

PRINCIPLE 2: ISOLATION

- Communicate only by message passing
- A fault in one process should be contained
- Share nothing concurrency
 Also known as the Actor model

wheessage passing should be contained rency Actor model

PRINCIPLE 3: LET IT CRASH

- Don't make things worse

- Will never be able to eliminate all faults Instead: plan for faults to happen • If a process hits an error, just crash it • Let someone other process fix/restart

THE ERLANG LANGUAGE

- Designed for telecom applications Soft real-time, highly reliable
 - Can swap in code updates live
- At the core: processes, messages, isolation

• Designed for processes that live forever

BIGIMPACT

• Runs Ericsson telecom switches Handles estimated 50% of all cell traffic OTP libraries, Open Telecom Protocol Runs Whatsapp and FB chat (previously) Whatsapp: 50 employees for 900M users (2015) Many successful applications CouchDB, Riak, Elixir, ...

SPAWNING PROCESSES

my_proc = fun() -> 2 + 2 end.
p_id = spawn(my_proc).

• Just like in Rust: pass it a closure

SENDING MESSAGES

```
p id ! hello.
self() ! there.
```

 Asynchronous channels: send never blocks Send directly to process, not to specific channel

RECEIVING MESSAGES

```
dolphin() ->
  receive
    do a flip ->
      io:format("How about no?~n");
    fish ->
      io:format("So long and thanks for all the fish!~n");
      ->
      io:format("Heh, we're smarter than you humans.~n")
  end.
```

• Do a case analysis on received message

• Each process has one incoming queue, like a mailbox

JOE'S FAVORITE PROGRAM

• Universal Server: can turn into any another process

universal server() -> receive {become, New proc} -> New proc() end.

• Lowercase match on string, uppercase variable

FACTORIAL SERVER • Wait for message, respond with factorial

factorial server() -> receive $\{From, N\} \rightarrow$ From ! factorial(N), factorial server() end.

factorial(0) $\rightarrow 1;$ factorial(N) \rightarrow N * factorial(N-1).

BECOMING FACTORIAL

• Turn a universal server into a factorial server

main() ->univ pid = spawn(fun universal server/0), univ pid ! {become, fun factorial server/0}, univ pid ! {self(), 50}, receive Response -> Response end.

/ 0 means zero arguments (Erlang dynamically typed)