LECTURE 22

Theory and Design of PL (CS 538) April 13, 2020



SHARED-STATE CONCURRENCY

IDEA: IT'S GOOD TO SHARE

- All threads can modify shared data
- Benefits
 - Threads can work on data in place
 - Big savings if the shared data is big
 - One, consistent view of shared data
- Drawbacks

 - Need to prevent data races

Not always safe to interrupt threads

CRITICAL SECTIONS

- All parts of code accessing shared piece of data
- Only one thread allowed in section at a time
 Otherwise: possible race condition
- Other threads must wait ("block") until safe to enter

shared piece of data n section at a time condition 'block") until safe to enter

A BANK ACCOUNT

 Bank account tracks current account balance Operations to deposit/withdrawn money

struct Account { balance: i32 }; impl Account { fn deposit(&mut self, amt: i32) { /* ... */ }

fn withdraw(&mut self, amt: i32) { if self.balance < amt {</pre> println!("Insufficient funds."); } **else** { self.balance -= amt;)

WHAT COULD GO WRONG?

- Suppose acct: Account has balance 100
- Suppose two threads call acct.withdraw(75)
- Interleaving may cause negative balance 1. T1 checks if balance is enough: OK 2. Before deduct balance, T2 runs 3. T2 checks if balance is enough: OK 4. T1 deducts 75: balance is 25 5. T2 deducts 75: balance is -25

LOCKS (MUTEXES)

- Threads can acquire or release locks If T1 tries to acquire lock held by T2, it has to wait until the lock is released by T2
- Main tool to enforce critical sections: locks • Only one thread can hold a lock at a time
- Prevent data races by carefully using locks

A BANK ACCOUNT, BUT SAFER

Following code gives the idea Note: not real Rust code (stay tuned)

struct Account { balance: i32, mutex: Mutex };

```
impl Account {
    fn withdraw(&mut self, amt: i32) {
        self.mutex.get lock();
        if self.balance < amt {</pre>
            println!("Insufficient funds.");
        } else {
            self.balance -= amt;
        self.mutex.unlock();
```

WHY DOES THIS WORK?

- Only one thread can get the lock
- Programmer doesn't pick which thread goes first
- If other thread tries to get lock, it must wait (block) • Result: thread runs withdraw with no interruptions

LOCKING DISCIPLINE

- More locks = more problems
- Programmer must coordinate how threads use locks
 - Which locks to acquire, when and where
 - What order to acquire locks
 - When and where to release locks
- Specified by programmer, not checked by compiler

EXAMPLES

- "When reading/writing, must hold global lock"
 Strongly limits concurrency
- "When reading/writing x, must hold lock for x"
 What if you need operate on two variables?
- Real schemes are usually much more complicated

MANY BUGS ARE POSSIBLE

- Too little locking or forget to take lock?
 Data races and unpredictable behavior
- Too much locking?
 Lots of threads wait
- Forget to release lock?
 Waiting threads will block forever
- Lots of threads waiting, little concurrency Forget to release lock?

MESSAGE-PASSING CONCURRENCY

IDEA: SHARING IS A BAD IDEA

- Don't share data between threads Each thread operates on private data only Threads send/receive messages
- Benefits
 - All interaction confined to thread input queues No sharing = no data races = no manual locks
- Drawbacks
 - Inefficient if we need to send lots of data

Harder to synchronize, no common view of data

A BANK ACCOUNT, MESSAGE PASSING

Idea: Clients send withdrawl messages
Fancier: Clients can have message queues too

struct Client { the_acct: &mut Account };
impl Client {
 fn withdraw(&self, amt: i32) {
 the_acct.send_withdraw(amt);
 }
}

A BANK ACCOUNT, MESSAGE PASSING Idea: give Account a message queue Again: not real Rust code (stay tuned)

enum AcctMsg { Withdraw(i32), // ... other messages ... } struct Account { balance: i32, msg: VecDeque<AcctMsg> };

```
impl Account {
  fn send withdraw(&mut self, amt: i32) {
    self.msg.push back(Withdraw(amt));
  fn run acct(&mut self) {
    loop \{
      if let Some(Withdraw(amt)) = self.msg.pop front() {
       if self.balance < amt {</pre>
          println!("Insufficient funds.");
        } else { self.balance -= amt; }
```

WHAT'S THE POINT?

- Single thread: no overlapping withdraws! If two clients send msgs, update queue correctly

- Account and Clients run in separate threads • Account processes messages one at a time Synchronization needed only at message queue • Reduce and restrict shared state
- As much as possible: run logic in a single thread

CHANNELS

- Central abstraction for message passing
- Goes from thread(s) to other thread(s)
- Messages might arrive in any order
- Receiver sees a single stream of messages

or message passing to other thread(s) re in any order e stream of messages

CHANNEL OPERATIONS

- Each thread can send or receive message via channel
- Synchronous channels
 - Receiving waits until something arrives (blocking)
- Asynchronous channels
 - Try-receive: does not block if nothing incoming
 - Select: wait for msg. from any of these channels

DEADLOCK



CIRCULAR WAIT

- All threads blocked waiting for other threads
- Shared-state example
 - T1: take lock x, take lock y
 - T2: take lock y, take lock x
- T1 takes x and T2 takes y: Deadlock! Two threads both waiting for message
- Obvious here, but harder with more locks/threads... • Can happen under message-passing too

DINING PHILOSOPHERS

- N philosophers (threads), sitting in a circle
- N forks (locks), one between every 2 philosophers
- Philosophers think, then eat, then think, ...
- To eat: philosopher takes left fork, then right fork
- If all philosophers take left fork: stuck!

, sitting in a circle een every 2 philosophers eat, then think, ... left fork, then right fork ft fork: stuck!

HOW TO FIX?

- One special philosopher: take right fork, then left
- Other philosophers: take left fork, then right
- Break the symmetry between threads
- In general: fixing deadlocks is very challenging

take **right** fork, then left left fork, then right ween threads ks is very challenging

RUST THREADS

SPAWNING THREADS

• Rust function: thread::spawn Caller passes in a closure for new thread to run • Terminology: caller is parent, new thread is child

```
fn main() {
  thread::spawn( || { // begin closure (child runs this)
      for i in 1..10 {
       // do some stuff, sleep for 10 seconds, ...
         // end closure
```

println!("hi number {} from the spawned thread!", i);

RETURNS IMMEDIATELY

• Returns a handle, parent continues running Child thread may not finish before parent

```
fn main() {
  let child = thread::spawn( || { // start child
      for i in 1..10 {
        // do some stuff, sleep for 10 seconds, ...
      } });
```

// parent thread continues for i in 1..5 {

println!("hi number {} from the main thread!", i); // do some stuff, sleep for 10 seconds, ...

println!("hi number {} from the spawned thread!", i);

JOINING THREADS

• join: wait for threads to finish Call with handle from spawn to wait for that thread

fn main() { let child = thread::spawn(|| { for i in 1..10 { // do some stuff, sleep for 10 seconds, ... } });

for i in 1..5 { println!("hi number {} from the main thread!", i); // do some stuff, sleep for 10 seconds, ...

child.join(); // wait for child to finish

println!("hi number {} from the spawned thread!", i);

THREAD ENVIRONMENT

Must move or clone environment into each thread

 Common case: parent wants to put data into thread • Mechanism: closure mentions external variables

let env var: String = String::from("foo"); let child one = thread::spawn(|| { // Not OK: env var doesn't live long enough let my var = env var; println!("Child 1 printing: {}", my var); });

let child two = thread::spawn(move || { // OK: thread takes ownership of env var let my var = env var; println!("Child 2 printing: {}", my var); });

EACH DATA HAS ONE OWNER

Compiler: variable moved into at most one thread

let mut mut var = String::new("foo"); let child one = thread::spawn(move || { // OK: thread takes ownership of mut var mut var.push str(" and bar"); });

let child two = thread::spawn(move || { // Not OK: can't move mut var again! mut var.push str(" and baz"); });

Rust compiler prevents data races!

OTHER RACES POSSIBLE

- Many resources not covered by aliasing rules Printing lines to screen
 - Reading/writing from file system

 - Sending/receiving network packets
- Races in other effects not controlled by Rust

WHAT ABOUT SHARING?

- Can't share mutable access to same data
- But what about sharing immutable access?
- Does this work?

let env var = String::new("foo"); let child one = thread::spawn(|| { // Thread borrows env var immutably println!("Child 1 says: {}", env var); });

let child two = thread::spawn(|| { // Thread borrows env var immutably println!("Child 2 says: {}", env var); });

DOESN'T LIVE LONG ENOUGH

- Parent may finish early, deallocate env var
- Child threads may hold dangling reference...
- Try: use Rc to share ownership

let env var old = Rc::new(String::new("foo")); let env var one = Rc::clone(&env var old); let env var two = Rc::clone(&env var old); let child one = thread::spawn(move || { println!("Child 1 says: {}", env var one); });

let child two = thread::spawn(move || { println!("Child 2 says: {}", env var two); });

// String will live as long as one Rc is still alive

STILL NOT HAPPY

 "Rc doesn't implement Sync, try Arc" Now it works. But what was the problem?

let env var old = Arc::new(String::new("foo")); let env var one = Arc::clone(&env var old); let env var two = Arc::clone(&env var old); let child one = thread::spawn(move || { println!("Child 1 says: {}", env var one); });

let child two = thread::spawn(move println!("Child 2 says: {}", env

TRAITS FOR THREAD-



THREAD-SAFETY

- Thread safe data: safe to share between threads
- Interface and internals must be carefully designed
 - Multiple threads may operate on same data
 - Threads may call different operations
 - Resumed/paused/interrupted at any time

hare between threads ust be carefully designed perate on same data ent operations rupted at any time

CHECKED BY RUST COMPILER

By default, custom types are not thread safe
If you share between threads, bad stuff happens
Much of standard library is thread safe
Vec, HashMap, String, ...

Rust compiler complains if you don't use thread safe libraries with threads!

THREAD-SAFETY TRAITS

- Tracked at the type level through traits Send trait: can be sent to another thread Sync trait: can be shared by multiple threads • Marker traits: no required implementations Can't implement in safe Rust, essentially a promise

- Examples

 - Rc doesn't implement Send or Sync: not safe! Arc implements Send and Sync: thread safe!